

DEDUCTIVE SYSTEMS WITH UNIFIED MULTIPLE-CONCLUSION RULES

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Following the approach suggested by Brentano and accepted and developed by Łukasiewicz, we study the deductive systems that treat asserted and rejected propositions equally, in the same way.

By statement, we understand the expressions of form $\oplus A$ - "A being asserted", and $\ominus A$ - "A being rejected". Accordingly, by a consequence relation we understand a consequence relation between sets of statements and statements.

Unified deductive systems are the ordered pairs (Γ, \mathbf{R}) , where Γ is a set of axioms (statements of form $\oplus A$) and anti-axioms (statements of form $\ominus A$), while \mathbf{R} is a set of multiple-conclusion rules – rules of form Δ/Ξ , where Δ and Ξ are the finite, perhaps empty, sets of statements. \blacktriangledown denotes the empty set of premises, and \blacktriangle denotes the empty set of conclusions.

Logic is coherent if the rule $\mathbf{Co} := \oplus p, \ominus p/\blacktriangle$ is admissible in it; and logic is full if the rule $\mathbf{Fu} := \blacktriangledown/\oplus p, \ominus p$ is admissible in it. The rules \mathbf{Co} and \mathbf{Fu} can be viewed as the rules expressing the law of non-contradiction and the law of the excluded middle.

Inclusion of the rules \mathbf{Co} and \mathbf{Fu} in a unified deductive system makes the system more flexible. In particular, adding these rules and the rule $\oplus \neg p/\ominus p$ to the classical propositional calculus (with postulated rule of substitution), results a unified deductive system that provides a complete Łukasiewicz style axiomatization: statement $\oplus A$ is derivable if and only if formula A is a classical theorem, while statement $\ominus A$ is derivable if and only if formula A is not a classical theorem.

Use of rules \mathbf{Co} and \mathbf{Fu} in formal derivations allows us to reduce multiple-conclusion rules to single-conclusion ones.

A unified deductive system that provides a complete and sound axiomatization of the classical logic with refutations is presented.